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## EUROPEAN PATENT APPLICATION

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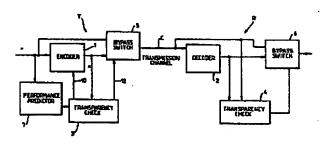
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- Data compression.
- A method of processing data for transmission from a transmitter T to a receiver R both of which are switchable between a compression mode in which the data stream is encoded or decoded with a data compression algorithm and a transparent mode in which the data stream is transmitted or received unencoded, comprising reading an input data stream at the transmitter T, encoding at least part of the input data stream with the data compression algorithm to form a compressed data stream, montoring the efficiency of compression of the compressed data stream and controlling the switching of the mode of the transmitter T and receiver R so that the input data stream is transmitted efficiently.

In the transparent mode both the receiver and the transmitter check independently the efficiency of compression and switch to the compression mode if the transparent mode is determined to be inefficient. In the compression mode the transmitter T determines the efficiency of compression and transmits a control code to the receiver R, switching both the receiver and transmitter to the transparent mode if the compression mode is determined to be inefficient.



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The present invention relates to data compression systems. Most useful forms of data have a significant degree of redundancy and as a result it is often advantageous to use a data compression system to encode the data in such a way that - redundant signals are eliminated, thereby reducing the space needed to store the data or reducing the bandwidth required to transmit the data.

It is known to use a method of processing data for transmission from a transmitter to a receiver both of which are switchable between a compression mode in which data stream is encoded or decoded with a data compression, algorithm and a transparent mode in which the data stream is transmitted or received unencoded, comprising reading an input data stream at the transmitter, encoding at least part of the input data stream with the data compression algorithm to form a compressed data stream, monitoring the efficiency of compression of the compressed data stream and controlling the switching of the mode of the transmitter and receiver so that the input data stream is transmitted efficiently. Car North

There is a significant overhead associated with the use of data compression algorithms and as a result in those parts of an input data stream which have a low degree of redundancy the use of the compression algorithm can result in a compression ratio less than unity, that is the "compressed" data occupies more space than the uncompressed data. It has been proposed to use methods of compression such as that described above in which the efficiency of the compression process is monitored and the transmitter and receiver are switched between a compression mode in which the data is encoded using a compression algorithm and a transparent mode in which the data is transmitted without compression as appropriate to ensure maximum efficiency. In one such system the contents of an entire data buffer are encoded using a data compression algorithm and the resulting data examined to determine whether it does in fact occupy less space. The data from the shorter of the compressed/uncompressed buffers is then transmitted together with a header which indicates whether the transmitted data is compressed or not. The decoder at the receiver then operates in the transparent or the compressed mode according to the value of the header. This scheme has the disadvantage that it is inflexible since it can only operate on complete buffers and is inefficient when there are are significant variations in the nature of the data within a single buffer. The need to encode an entire buffer before transmission introduces a significant delay and moreover there is a significant

normoverhead associated with all data since uncompreswied data still needs to be prefixed with a code to Indicate that the receiver should operate in the transparent mode. As an alternative to accumulating an entire buffer before determining whether the is compressed or stransparent mode is appropriate and schemes have also been oproposed in which a sample from the beginning of the data stream is tested to determine which method of transmission how is most refficient and the rest of the data stream as transmitted; and received in the compression or transparent mode according to the results of that determination on the sample. This implementation of the method outlined above suffers the disadvan-157 / tage that the overall efficiency of transmission can god be seriously effected if the initial sample from the 。接頭はdata istreamais lunrepresentative of the data as a 。就是接到1.00mm的方面。

we build According to the present invention, in such a 20/38/method thithe transparent mode both the receiver glassy and the transmitter check independently the effibe able compression, and switch to the compresand existing mode if the transparent mode is determined to when be inefficient, and in the compression mode the 25 transmitter determines the efficiency of compression and transmits a control code to the receiver. both the receiver and transmitter thereby switching to the transparent mode if the compression mode The largest determined to be inefficient. The largest largest at

90,568 transmission of additional characters from had the atransmittor to the receiver to provide control codes which switch the receiver between modes has a serious effect on the efficiency of the trans-, a, mission, system when the system is operating in 35: In the transparent mode, in the compression mode by contrast the use of the compression algorithm frees bandwidth on the transmission channel which can be used for the control codes. The present inventor has found that significant advantages arise from the 40 is method of the present invention, in which when in the transparent mode the transmitter and receiver monitor the efficiency of compression independently and so switch between modes where appropriate and without requiring the transmission of control codes. 45% public when the system is in the compression mode the transmitter transmits a code to cause the recalciver to switch between modes so that the state of and the tracelyer is dependent upon the state of the transmitter. The provision of dependent switching 50 when the system is in the compression mode enables the method of the present invention to take advantage of the fact that the transmitter, unlike the receiver, has access to data which has not yet been transmitted and can look ahead to detect changes in the data characteristics and where ap-

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propriate elect to switch to the transparent mode before the performance of the system deteriorates. Prompt switching between modes has found to be critical in the compression mode where a compression ratio less than unity can result in a severe degradation in performance in the transparent et et mode by contrast, the compression ratio is fixed at with unity providing a lower limit on the refficiency of graphe systematic section is at a specific to the con-

some a Preferably sine the transparent mode both the receiver and transmitter determine independently the efficiency of compression of current data and in the compression mode the transmitter determines the efficiency of compression of untransmitted data in the data stream. Preferably the transmitter and ... receiver determine the efficiency of compression of the current data by comparing the size of a seequence, of data-symbols: from the current data before and after compression and the transmitter determines the efficiency of compression of untransmitted data from the data stream by calculating from the untransmitted data: a parameter indicative of its compressibility.

The data compression algorithm employed by the receiver and transmitter in the compression mode may be non-adaptive so that the code words allocated to given input symbols are predetermined but preferably the data compression algorithm is adaptive, the code words allocated to input symbols being modified in accordance with the char-. a ...acteristics/of the data stream... ................... below at le

A system in accordance with the present invenet ation will mow beadescrized in detall with reference to the figure of the accompanying drawings which is a block diagram showing a receiver and transmilitar suitable for use in the method of the present a rad**invention**es, levicos, en local la fosse efecto

A data transmission system comprises a transat amitter T linked via attransmission channel C to a receiver Rathe transmission channels may be provided by an GSTN finking remote stations housing the transmitter Trand receiver R respectively or may comprise a mass storage device to which the transmitter: Towrites data and from which the rea ceiver R reads data. In this latter case the transmitter Trand receiver R are formed in a single device.

The transmitter T includes an encoder 1 and a This in a compression mode data from the input is fed to the encoder 1 where it is encoded using a data compression algorithm. In the present example the algorithm used is the well known Huffman code described in the paper " A Method To?" the Construction of Minimum Redundancy Codes" by D.A. Huffman, proceedings of the IRE, Vol. 40,8 1952. A sind decoder 2 in the receiver R decodes the received 🌃 data using a complementary algorithm.

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the transmitter and receiver operate to bypass the encoder and decoder so that data is transmitted and received in uncompressed form. Both the transmitter T and receiver R include transparency checking mechanisms 3,4 which, in the manner more fully described below, determine the efficlency of compression of the current data-handled by the encoder and decoder. The transmitter includes in addition a performance predictor 7 which 10: (analyses data in advance of its encoding. The performance predictor 7 provides output for the transparency check mechanism 3 to cause the system to switch between modes when a change in the nature of the data is detected.

In use, the system starts in transparent mode, the transmitter bypass switch 5 selecting the input data and the decoder bypass switch accepting Sunencoded date from the transmission channel T. The transparency checking mechanisms 3, 4, in 20 both the receiver and transmitter lavout independently the following procedure. درية الأرتقاب

(i) Identify a known starting point, e.g. the 2.60 start of the message and declared above the

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The number of bits transmitted some starting point was a second to the

34 5 34 (fill) Calculate the number of bits that would and been transmitted if compression had been applied

(Iv) At predetermined intervals compare the so we results of steps (ii) and (iii) was a same

(v) If the number of bits from (III) is less than the number of bits from (II) switch to compression mode and begin transmitting/receiving the date in compressed form.

33 in compression mode the encoder compresses the data stream prior to transmission and the decoder performs the reverse function. The transparency checking mechanism in the receiver carries Cout the following steps: See the manufactor will be a see that

(I) Identify a known starting point

(ii) Count the number of bits transmitted

\*\* \*\* (iii) Calculate the number of bits that would have been transmitted if compression had not been 300 V. applied

48 12 42 F (iv) At predetermined intervals compare the results of step (ii) with the results of step (iii)

(v) if the number of bits from (iii) is less than the number of bits from (ii) transmit an explicit code word to the decoder to signal transition to transparent mode; and begin transmitting the data likar i enu s in uncompressed form.

The performance predictor 7 in the transmitter carries out the following steps when the system is 2.4 in the compression mode:

66 -(i) Examine the data as yet untransmitted for characteristics that indicate that the performance of the data compression encoder may be degraded

(ii) If the test applied in (i) Indicates that the . And this of the contract

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performance would be degraded if the encoder remains in compressed mode transmit an explicit code word to the decoder to signal transition to transparent mode and begin transmitting the data in uncompressed form.

In compression mode the decoder:

- (i) receives and decodes code words
- (ii) when an explicit code word is received indicating that a transition to transparent mode is being made switches to transparent mode.

The system outlined above provides improved efficiency in both transparent and compressed modes. In compressed mode the encoder looks ahead for changes in the data characteristics and can elect to switch to transparent mode before performance deteriorates. In transparent mode no transmitted bandwidth is wasted on explicit signals transmitted between the transmitter and receiver to indicate the change of state.

Alternative arrangements are possible in which both transmitter and receiver rely solely upon checks carried out on current data without use of a predictive function.

As outlined above the transparency check function 3 in the transmitter counts the number of bits entering the encoder 1 and the number of bits leaving the encoder 1. After a pre-defined interval, which in the present example comprises 256 input characters, the two counts are compared and if the number of bits leaving the encoder is less than the number of bits entering it then the bypass switch 5 is switched by means of a control lead 12 to accept input from the encoder 1 thereby from that instant onwards transmitting the compressed data stream over the transmission channel.

The transparency checking mechanism 4 in the receiver performs an equivalent function to that in the transmitter in transparent mode and may optionally do so as well in compression mode. The transparency checking mechanism 4 includes a copy of the encoder function which it uses to calculate the same count of encoded and unencoded bits as the corresponding transparency checking mechanism 3 in the encoder. The two transparency checking mechanisms 3,4 employ identical logical decision functions and so switch from transparent to compressed modes at identical points within the transmitted bit stream.

The performance predictor 7 in the transmitter may take a number of different forms, but in the present example counts the frequency of occurrence of the characters within the input data stream in advance of their being encoded and calculates the following prediction function which produces an output dependent on the average length of the characters:

(i) Let the input characters be drawn from the set c[0].c[n] and let the character size be S

bits (typically 7 or 8 bits)

the folice of the process of a filter folices broaders.

(ii) Let the counted occurrence frequency for each character be f[0]..f[n]

(iii) Let the codeword length used within the encoder (1) for each character be L[0]..L[n] bits

(iv) Let Sum (x[i]) represent the sum for i=0 to n of x[i]

(v) Calculate A = Sum(f[i].L[i])/Sum(f[i])

If the value of A is greater than the number of input characters then the encoder performance would deteriorate and the transparency check function 3 signals via a control lead 10 that the encoder 3 should emit a special control codeword. After emitting the special control codeword the bypass switch 5 is switched to accept uncompressed input data and the transmitter enters the transparent state.

On receiving the special control codeword the receiver discards the codeword, switches the bypass switch 6 to accept undecoded input data and enters the transparent state.

Although the scheme is described above with reference to the Huffman algorithm which is non-adaptive the present invention may equally be used with adaptive compression systems in which the codewords assigned in the compression mode are varied to suit the current characteristics of the data stream. In this case the predictive function is varied appropriately to match the encoding function used.

## Claims

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1. A method of processing data for transmission from a transmitter to a receiver both of which are switchable between a compression mode in which a data stream is encoded or decoded with a data compression algorithm and a transparent mode in which the data stream is transmitted or received unencoded, comprising reading an input data stream at the transmitter, encoding at least part of the input data stream with the data compression algorithm to form a compressed data stream, monitoring the efficiency of compression of the compressed data stream and controlling the switching of the mode of the transmitter and receiver so that the input data stream is transmitted efficiently, in which in the transparent mode both the receiver and the transmitter check independently the efficiency of compression and switch to the compression mode if the transparent mode is determined to be inefficient, and in the compression mode the transmitter determines the efficiency of compression and transmits a control code to the receiver, both the receiver and transmitter thereby switching to the transparent mode if the compression mode is determined to be inefficient.

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2. A method according to claim 1, in which in the compression mode the transmitter monitors the efficiency of compression of untransmitted data in f . . . . . .

the data stream.

3. A method according to claim 2, in which the transmitter monitors the efficiency of compression of untransmitted data in the data stream by calculating from the untransmitted data a parameter Indicative of its compressibility.

4. A method according to any one of the preceding claims, in which the transmitter and receiver monitor the efficiency of compression of current data by comparing the size of a sequence of data , symbols from the current data before and after compression.

5. A method according to any one of the preceding, claims, in which the data compression algorithm is non-adaptive so that the code words allocated to given input symbols are predetermined. Check to at the price to the

8. A method according to any one of claims 1 to 4, in which the data compression algorithm is adaptive, the code words allocated to input symbols being modified in accordance with the characteristics of the data stream.

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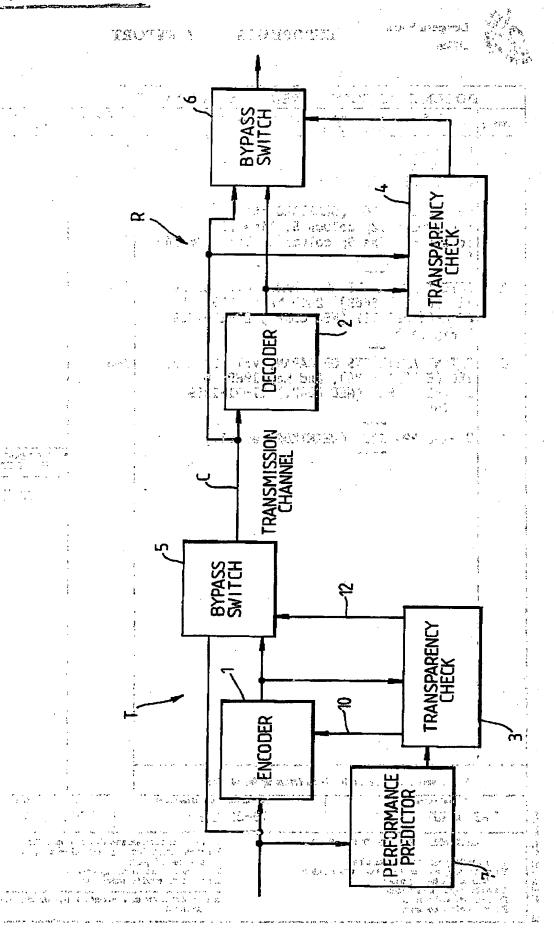
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## **EUROPEAN SEARCH REPORT**

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EP 89 31 2782

ategory	Citation of document with indication, where appropriate, of relevant passages				CLASSIFICATION OF THE APPLICATION (Int. Cl.5)
X	* Figures	9 038 (HIR/ 1-14; colum , line 10 *	1	H 03 M 7/48	
X come vigora	US-A-3 185 824 (BLASBALG et al.)  * Figures 1,2; column 5, line 70 - column 6, line 9; column 6, lines 46-74  *  PATENT ABSTRACTS OF JAPAN, vol. 13, no. 121 (E-732)[3469], 24th March 1989; & JP-A-63 290 021 (NEC CORP.) 28-11-1988  * Figure *				
<b>A</b>					
E	PATENT AB 186 (E-75 JP-A-1 12 * Figure	2)[3534], 21 621 (NEC CO	JAPAN, vol. 13, no. nd May 1989; & DRP.) 17-01-1989	1-4	
A	US-A-3 39	4 352 (WERI	VIKOFF et al.)		TECHNICAL FIELDS SEARCHED (Int. Cl.5)
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- X: particularly relevant if taken alone
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  A: technological background
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  P: intermediate document

- E: earlier patent document, but publisher the filing date

  D: document cited in the application

  L: document cited for other reasons

- & : member of the same patent family, corresponding document

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